

Positive and completely positive maps via free additive powers of probability measures

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St John's, January 25th, 2013

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Entanglement in Quantum Information Theory

- Quantum states with n degrees of freedom are described by **density matrices**

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- Equivalently, $\mathcal{SEP} = \text{conv} [\mathbb{M}_m^{1,+} \otimes \mathbb{M}_n^{1,+}]$.
- Non-separable states are called **entangled**.

More on entanglement - pure states

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- For rank one quantum states, entanglement can be detected and quantified by the **entropy of entanglement**

$$E_{\text{ent}}(P_x) = H(s(x)) = - \sum_{i=1}^{\min(m,n)} s_i(x) \log s_i(x),$$

where $x \in \mathbb{C}^m \otimes \mathbb{C}^n \cong \mathbb{M}_{m \times n}(\mathbb{C})$ is seen as a $m \times n$ matrix and $s_i(x)$ are its singular values.

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- A pure state $x \in \mathbb{C}^m \otimes \mathbb{C}^n$ is separable $\iff E_{\text{ent}}(P_x) = 0$.

An image of entanglement



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Shor's quantum factoring algorithm

- Runs on a quantum computer with polynomial time $O(\log^3 N)$.
- Classical sieve algorithms run in sub-exponential time $O(\exp(\log^{1/3} N))$.
- Entanglement is necessary for the exponential speed-up.
- State of the art factorization in labs: $21 = 3 \times 7$ [2011], 143 (?) [2012].

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- Detecting entanglement for general states in low dimension $\mathbb{C}^2 \otimes \mathbb{C}^2$ and $\mathbb{C}^2 \otimes \mathbb{C}^3$ is possible via the **PPT criterion** [Horodecki].

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- Detecting entanglement for general states in low dimension $\mathbb{C}^2 \otimes \mathbb{C}^2$ and $\mathbb{C}^2 \otimes \mathbb{C}^3$ is possible via the **PPT criterion** [Horodecki].
- In general, there exists a countable hierarchy of conditions characterizing separability [Doherty et al] that can be checked by semidefinite programming.

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- Let $\mathcal{N} : \mathbb{M}_n \rightarrow \mathcal{A}$ be a **completely positive** map. Then, for **every** state $\rho_{12} \in \mathbb{M}_{mn}^{1,+}$, one has $[\text{id}_m \otimes \mathcal{N}](\rho_{12}) \geq 0$.

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- Hence, positive, but not CP maps \mathcal{N} provide **sufficient entanglement criteria**: if $[\text{id}_m \otimes \mathcal{N}](\rho_{12}) \not\geq 0$, then ρ_{12} is entangled.

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- Hence, positive, but not CP maps \mathcal{N} provide **sufficient entanglement criteria**: if $[\text{id}_m \otimes \mathcal{N}](\rho_{12}) \not\geq 0$, then ρ_{12} is entangled.
- Moreover, if $[\text{id}_m \otimes \mathcal{N}](\rho_{12}) \geq 0$ for **all** positive, but not CP maps \mathcal{N} , then ρ_{12} is separable.

Positive Partial Transpose matrices

- The transposition map t is positive, but not CP. Define the convex set

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- For $(m, n) \in \{(2, 2), (2, 3)\}$ we have $\mathcal{SEP} = \mathcal{PPT}$. In other dimensions, the inclusion $\mathcal{SEP} \subset \mathcal{PPT}$ is strict.

The PPT criterion at work

- Recall the Bell state $\rho_{12} = P_{\text{Bell}}$, where

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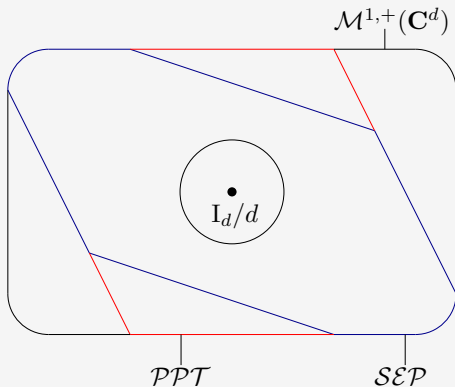
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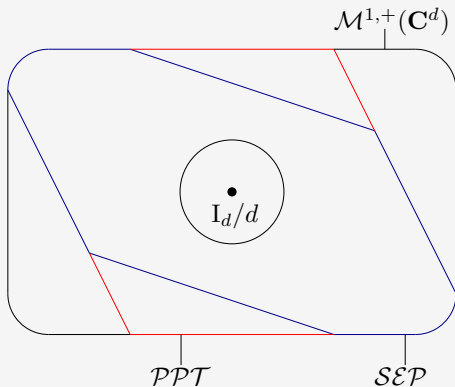
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- This matrix is no longer positive \implies the state is entangled.

Three convex sets

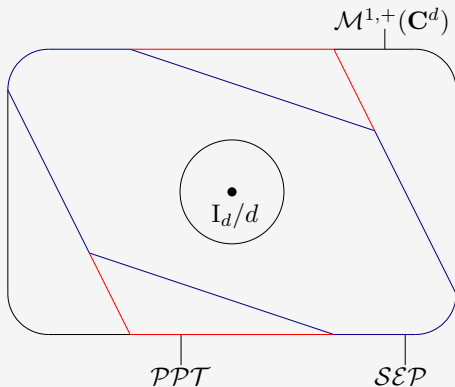


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- All these sets contain an open ball around the identity.

The Choi matrix of a map

- For any n , recall that the **maximally entangled state** is the orthogonal projection onto

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Theorem (Choi '72)

A map $\mathcal{N} : \mathbb{M}_n \rightarrow \mathcal{A}$ is CP *iff* its Choi matrix $C_{\mathcal{N}}$ is positive.

The Choi-Jamiolkowski isomorphism

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- The map $\mathcal{N} \mapsto C_{\mathcal{N}}$ is called the **Choi-Jamiolkowski** isomorphism.
- It sends:
 - 1 All linear maps to all operators;
 - 2 Hermiticity preserving maps to hermitian operators;
 - 3 Entanglement breaking maps to separable quantum states;
 - 4 Unital maps to operators with unit left partial trace ($[\text{Tr} \otimes \text{id}]C_{\mathcal{N}} = I_d$);
 - 5 Trace preserving maps to operators with unit left partial trace ($[\text{id} \otimes \text{Tr}]C_{\mathcal{N}} = I_n$).

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Theorem

A map $\mathcal{N} : \mathbb{M}_n \rightarrow \mathcal{A}$ is k -positive iff its Choi matrix $C_{\mathcal{N}}$ is k -positive. In particular, \mathcal{N} is positive iff $C_{\mathcal{N}}$ is block-positive.

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Under the above assumptions, if $\text{supp}(\mu^{\boxplus n/k}) \subset (0, \infty)$ then, almost surely as $d \rightarrow \infty$, the map $\mathcal{N}_\mu^{(d)}$ is k -positive.

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- When $d \rightarrow \infty$, the spectrum of $A_d + B_d$ converges to $a \boxplus b$.

Proof ingredients

Let $\mathcal{N}_\mu^{(d)} : \mathbb{M}_n \rightarrow \mathbb{M}_d$ be the map whose Choi matrix is $U_d X_d U_d^*$.

Theorem

If $\text{supp}(\mu^{\boxplus n/k}) \subset (0, \infty)$ then, almost surely as $d \rightarrow \infty$, the map $\mathcal{N}_\mu^{(d)}$ is k -positive. If $\text{supp}(\mu^{\boxplus n/k}) \cap (-\infty, 0) \neq \emptyset$ then, almost surely as $d \rightarrow \infty$, $\mathcal{N}_\mu^{(d)}$ is not k -positive.

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Proposition

A map \mathcal{N} is k -positive iff for any self-adjoint projection $P \in \mathbb{M}_n$ of rank k , the operator $P \otimes I_{\mathcal{A}} \cdot C_{\mathcal{N}} \cdot P \otimes 1_{\mathcal{A}}$ is positive.

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Proposition (Nica and Speicher)

Let x, p be free elements in a ncps (M, τ) and assume that p is a selfadjoint projection of rank $\tau(p) = 1/t$ ($t \geq 1$) and that x has distribution μ . Then, the distribution of $t^{-1} p x p$ inside the contracted ncps $(p M p, \tau(p \cdot p))$ is $\mu^{\boxplus t}$

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Theorem

The map \mathcal{N}_μ is k -positive iff $\text{supp}(\mu^{\boxplus n/k}) \subseteq \mathbb{R}_+$.

Example: semicircular measures

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- We have $s_{a,\sigma}^{\boxplus n/k} = s_{an/k, \sigma\sqrt{n/k}}$, with support $\text{supp}(s_{a,\sigma}^{\boxplus n/k}) = [an/k - 2\sigma\sqrt{n/k}, an/k + 2\sigma\sqrt{n/k}]$.

Theorem

Let n be an integer and a, σ some positive parameters. The map $\mathcal{N}_{a,\sigma} : \mathbb{M}_n \rightarrow M$ associated to a semi-circular distribution $s_{a,\sigma}$ is k -positive iff $k \leq 4n\sigma/a^2$. In particular, for any n and any $k < n$, there exist parameters $a, \sigma > 0$ such that the above map is k -positive but not $k + 1$ -positive.

Merci !