

# Block-modified Wishart matrices and applications to entanglement theory

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*joint work with Teodor Banica (Cergy)*

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# Entanglement in Quantum Information Theory

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- ▶ Non-separable states are called **entangled**.

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- ▶ For rank one quantum states, entanglement can be detected and quantified by the von Neumann entropy

$$H(P_x) = S(\text{sv}(x)) = - \sum_i s_i(x) \log s_i(x), x \in \mathbb{C}^{d_1} \otimes \mathbb{C}^{d_2} \cong \mathcal{M}_{d_1 \times d_2}(\mathbb{C}).$$

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- ▶ Detecting entanglement for general states  $\mathbb{C}^2 \otimes \mathbb{C}^2$  and  $\mathbb{C}^2 \otimes \mathbb{C}^3$  is trivial via the **PPT criterion** [Horodecki].

## Positive Partial Transpose matrices

- ▶ A map  $f : \mathcal{M}(\mathbb{C}^d) \rightarrow \mathcal{M}(\mathbb{C}^d)$  is called
  - ▶ **positive** if  $A \geq 0 \implies f(A) \geq 0$ ;
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$$\mathbf{PPT} = \{\rho_{12} \in \mathcal{M}^{1,+}(\mathbb{C}^{d_1} \otimes \mathbb{C}^{d_2}) \mid [\text{id}_{d_1} \otimes t_{d_2}](\rho_{12}) \geq 0\}.$$



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- ▶ For  $(d_1, d_2) \in \{(2, 2), (2, 3)\}$  we have  $\mathcal{SEP} = \mathcal{PPT}$ . In other dimensions, the inclusion  $\mathcal{SEP} \subset \mathcal{PPT}$  is strict.

## The PPT criterion at work

- ▶ Recall the Bell state  $\rho_{12} = P_{\text{Bell}}$ , where

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- ▶ Written as a matrix in  $\mathcal{M}^{1,+}(\mathbb{C}^4)$

$$\rho_{12} = \frac{1}{2} \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{pmatrix} = \frac{1}{2} \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}.$$

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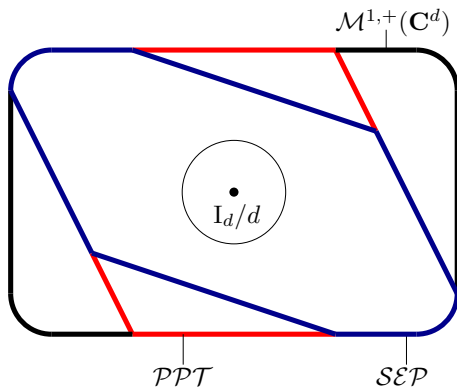
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- ▶ This matrix is no longer positive  $\implies$  the state is entangled.

## Three convex sets



- ▶ States in  $\mathcal{PPT} \setminus \mathcal{SEP}$  are called **bound entangled**: no “maximal” entangled can be distilled from them.
- ▶ All these sets contain an open ball around the identity.

## The problem we consider

$$\mathcal{M}^{1,+}(\mathbb{C}^{d_1 d_2}) = \{\rho \mid \text{Tr} \rho = 1 \text{ and } \rho \geq 0\}$$

$$\mathcal{SEP} = \left\{ \sum_i t_i \rho_1(i) \otimes \rho_2(i) \right\} = \text{conv} \left[ \mathcal{M}^{1,+}(\mathbb{C}^{d_1}) \otimes \mathcal{M}^{1,+}(\mathbb{C}^{d_2}) \right]$$

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### Problem

*Compare the convex sets*

$$\mathcal{SEP} \subset \mathcal{PPT} \subset \mathcal{M}^{1,+}(\mathbb{C}^{d_1 d_2}).$$

## Probability measures on $\mathcal{M}_d^{1,+}(\mathbb{C})$

- ▶ Let  $X \in \mathcal{M}_{d \times s}(\mathbb{C})$  a rectangular  $d \times s$  matrix with i.i.d. complex standard Gaussian entries. Define the random variables

$$W_{d,s} = XX^* \text{ and } \mathcal{M}_d^{1,+}(\mathbb{C}^d) \ni \rho_{d,s} = \frac{XX^*}{\text{Tr}(XX^*)} = \frac{W_{d,s}}{\text{Tr} W_{d,s}}.$$



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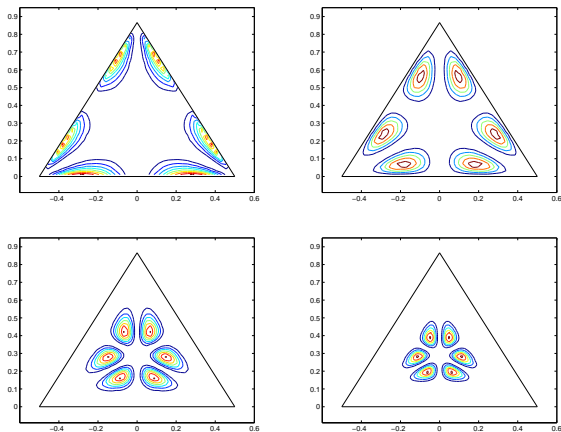
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- ▶ Almost surely,  $\rho_{d,s}$  has full rank iff  $s \geq d$ .
- ▶ The measure  $\mu_{d,s}$  is unitarily invariant: there exist a probability measure  $\nu_{d,s}$  on the probability simplex  $\Delta_d = \{\lambda \in \mathbb{R}^d \mid \lambda_i \geq 0, \sum \lambda_i = 1\}$  such that if  $\lambda \sim \nu_{d,s}$  and  $U$  is a Haar unitary matrix independent of  $\lambda$ ,

$$U \text{diag}(\lambda) U^* \sim \mu_{d,s}.$$

# Eigenvalues for induced measures



**Figure:** Induced measure eigenvalue distribution for  $(d = 3, s = 3)$ ,  $(d = 3, s = 5)$ ,  $(d = 3, s = 7)$  and  $(d = 3, s = 10)$ .

## Volume of convex sets under the induced measures

- ▶ Let  $C \subset \mathcal{M}^{1,+}(\mathbb{C}^d)$  a convex body, with  $I_d/d \in C^\circ$ . Then

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### Definition

A pair of functions  $s_0(d), s_1(d)$  are called a **threshold** for a family of convex sets  $\{C_d\}_{d \geq 2}$  if both conditions below hold

- ▶ If  $s(d) \lesssim s_0(d)$ , then

$$\lim_{d \rightarrow \infty} \mu_{d,s(d)}(C_d) = 0;$$

- ▶ If  $s(d) \gtrsim s_1(d)$ , then

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## Threshold for $\mathcal{SEP}$

Theorem (Aubrun, Szarek, Ye - 2011)

There exists a constant  $C$  such that the pair  $s_0 = Cd^{3/2}$ ,  $s_1 = Cd^{3/2} \log^2 d$  is a threshold for  $\mathcal{SEP}$ .

In other words, if  $s < Cd^{3/2}$ , then

$$\lim_{d \rightarrow \infty} \mu_{d,s}(\{\rho \text{ is entangled}\}) = 1$$

and if  $s > Cd^{3/2} \log^2 d$ , then

$$\lim_{d \rightarrow \infty} \mu_{d,s}(\{\rho \text{ is separable}\}) = 1.$$

# Partial transposition of a Wishart matrix

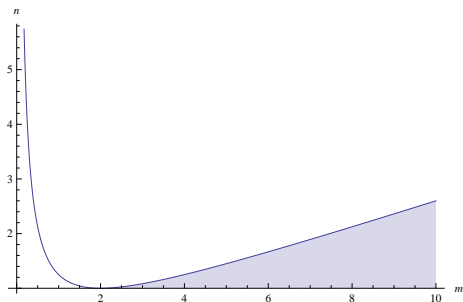
## Theorem (Banica, N.)

Let  $W$  be a complex Wishart matrix of parameters  $(dn, dm)$ . Then, with  $d \rightarrow \infty$ , the empirical spectral distribution of  $mW^\Gamma$  converges in moments to a **free difference of free Poisson distributions** of respective parameters  $m(n \pm 1)/2$ .

## Corollary

The limiting measure in the previous theorem has positive support iff

$$n \leq \frac{m}{4} + \frac{1}{m} \text{ and } m \geq 2.$$





## Threshold for $\mathcal{PPT}$ , unbalanced & balanced case

Theorem (unbalanced case, Banica, N.)

*In the unbalanced case  $d_1 = d \rightarrow \infty$ ,  $d_2 = n$  fixed, the lower bound of a threshold for  $\mathcal{PPT}$  is given by  $s_0 = \left[2 + 2\sqrt{1 - n^{-2}}\right] d$ .*

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### Theorem (balanced case, Aubrun - 2010)

*In the balanced case  $d_1 = d_2 = d \rightarrow \infty$ , a threshold pair for  $\mathcal{PPT}$  is given by  $s_0 = s_1 = 4d$ .*

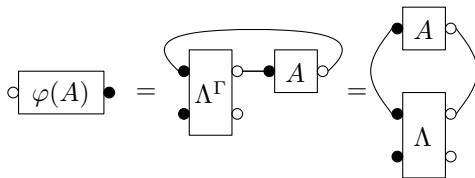
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- ▶ Define the **Choi matrix**  $\Lambda$  of  $\varphi$

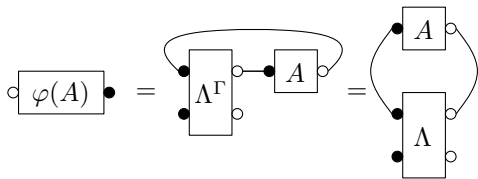
$$\varphi(A) = (\text{Tr} \otimes id)[(t \otimes id)\Lambda \cdot (A \otimes 1)]$$



## Generalizing partial transposition

- ▶ Replace the transposition map  $t$  with an arbitrary, **hermiticity preserving** linear map  $\varphi : \mathcal{M}(\mathbb{C}^d) \rightarrow \mathcal{M}(\mathbb{C}^d)$ .
- ▶ Define the **Choi matrix**  $\Lambda$  of  $\varphi$

$$\varphi(A) = (\text{Tr} \otimes \text{id})[(t \otimes \text{id})\Lambda \cdot (A \otimes 1)]$$



### Problem

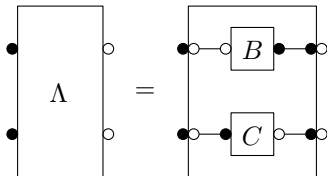
Compute the asymptotic spectrum of

$$\tilde{W} = (\text{id} \otimes \varphi)W,$$

where  $W$  is a Wishart random matrix,  $d \rightarrow \infty$  and  $n$  is fixed.

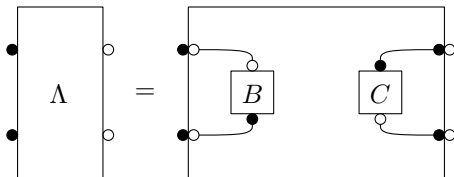
## Some examples

- ▶  $\varphi(A) = \text{Tr}(BA)C$ , in the case  $C = c1$ .
- ▶  $\Lambda = B^\top \otimes C$ .



## Some examples

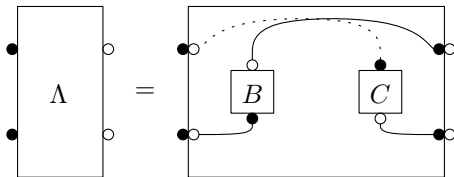
- ▶  $\varphi(A) = BAC$ , for any  $B, C$ .
- ▶  $\Lambda = |B\rangle\langle C|$ ,





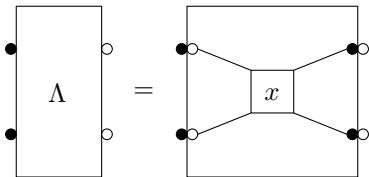
## Some examples

- ▶  $\varphi(A) = BA^tC$ , in the case  $BC = c1$ .
- ▶  $\Lambda = \text{SWAP}_{BC}$ ,



## Some examples

- ▶  $\varphi(A) = xA^\delta$ , in the case  $x = c1$ .
- ▶  $\Lambda = \text{Center}_x$ ,



## Our result

Theorem (Banica, N. - work in progress)

Let  $\tilde{W} = (id \otimes \varphi)W$ , where  $W$  is a complex Wishart matrix of parameters  $(dn, dm)$ , and where  $\varphi : M_n(\mathbb{C}) \rightarrow M_n(\mathbb{C})$  is a self-adjoint linear map, coming from a matrix  $\Lambda \in M_n(\mathbb{C}) \otimes M_n(\mathbb{C})$ . Then, under suitable “planar” assumptions on  $\varphi$ , we have  $\delta m \tilde{W} \sim \pi_{mn\rho} \boxtimes \nu$ , with  $\rho = \text{law}(\Lambda)$ ,  $\nu = \text{law}(D)$ ,  $\delta = \text{tr}(D)$ , where  $D = \varphi(1)$

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- ▶ Idea of the proof

$$\lim_{d \rightarrow \infty} (\mathbb{E} \circ \text{tr})((m\tilde{W})^p) = \sum_{\pi \in NC(p)} (mn)^{\#\pi} \text{tr}_{(\pi, \gamma)}(\Lambda).$$

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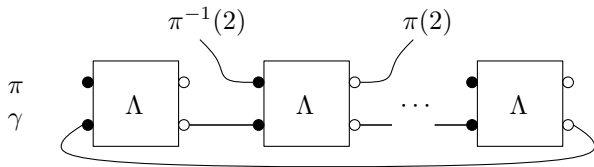
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- ▶ Identify the free cumulants, if the general term in the sum above is **multiplicative**.

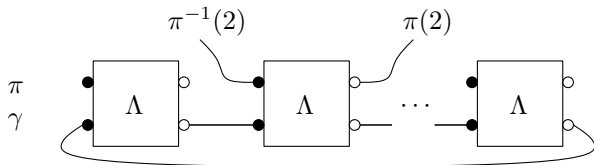
# Why does this fail for general $\varphi$ ?

- ▶ We have

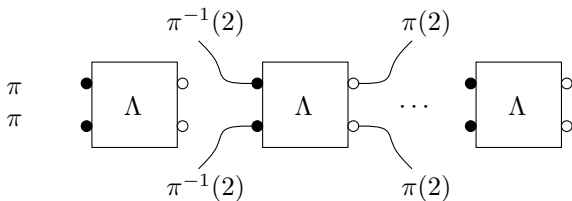


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# Thank you !

<http://arxiv.org/abs/1105.2556>

+

work in progress