

APPROXIMATING GS OF 1D GAPPED HAMIL. BY MPS (16/04/20)

→ Hilbert space $H = \bigotimes_{i=1}^N \mathbb{C}^d$; Hamiltonian $H = \sum_{i=1}^{N-1} H_{i,i+1}$ with $\|H_{i,i+1}\| \leq J$

→ We assume that H is **gapped**: $\text{spec } H = \{0\} \cup S$: $S \subseteq [\underbrace{\Delta E}_{\text{gap}}, +\infty)$

→ we say that $A \in \mathcal{B}(H)$ has support in X if $A = A_X \otimes I_{X^c}$

Theorem [LR bound '72] $\exists C, \xi, \nu > 0$ s.t. $\forall A, B \in \mathcal{B}(H)$ such that $\text{supp } A \subseteq X, \text{supp } B \subseteq Y$

$$\|[A, B(t)]\| \leq C \min(|X|, |Y|) \|A\| \|B\| \exp(2\xi^{-1}(\nu|t| - d(X, Y)))$$

Corollary $\forall |t| \leq \frac{d(X, Y)}{2\nu}$, $\|[A, B(t)]\| \leq C \min(|X|, |Y|) \|A\| \|B\| \exp\left(-\frac{d(X, Y)}{\xi}\right)$

→ Hastings' idea: use this to prove the area law for the GS of a gapped Hamiltonian

Definition $\psi_0 \in \mathcal{H}$, $\|\psi_0\| = 1$, $H\psi_0 = 0$ $\rho_{1,i} = \text{Tr}_{i+1, \dots, N} |\psi_0\rangle\langle\psi_0|$

Entanglement entropy $S(\rho_{1,i}) = -\text{Tr } \rho_{1,i} \log \rho_{1,i}$

→ trivial bound: $S(\rho_{1,i}) \leq i \log d$; when $S(\rho_{1,i}) \propto i \rightarrow$ **volume law**

when $S(\rho_{1,i}) = O(\log i) \rightarrow$ **area law**

Theorem [Hastings '07] $\exists S_{\max} > 0$ s.t. $\max_i S(\rho_{1,i}) \leq S_{\max}$ **Independent of N**

Actually, $S_{\max} = C_0 \xi' \ln \xi' \ln d \cdot 2^{\xi' \ln d}$ with $\xi' = 6 \min\left(\frac{2\nu}{\Delta E}, \xi\right)$

Today: use this entropy bound to $|\psi_0\rangle \approx \text{MPS}$ with "small" bond dimension

→ **Rényi entropies** $\alpha \geq 0, \alpha \neq 1$ $S_\alpha^\rho = \frac{1}{1-\alpha} \log \text{Tr}(\rho^\alpha)$

$$\rightarrow S_1(\rho) = -\text{Tr } \rho \log \rho = \lim_{\alpha \rightarrow 1} S_\alpha(\rho)$$

$$\rightarrow S_\infty(\rho) = -\log \|\rho\|_\infty = \log \lambda_{\max}(\rho)$$

$$\rightarrow S_0(\rho) = \log \dim \text{supp}(\rho)$$

→ **Matrix Product state** $(\mathbb{C}^d)^{\otimes N} \ni |\psi\rangle = \sum_{i_1, \dots, i_N} \text{Tr} [A_{i_1}^{[1]} \dots A_{i_N}^{[N]}] |i_1 \dots i_N\rangle$

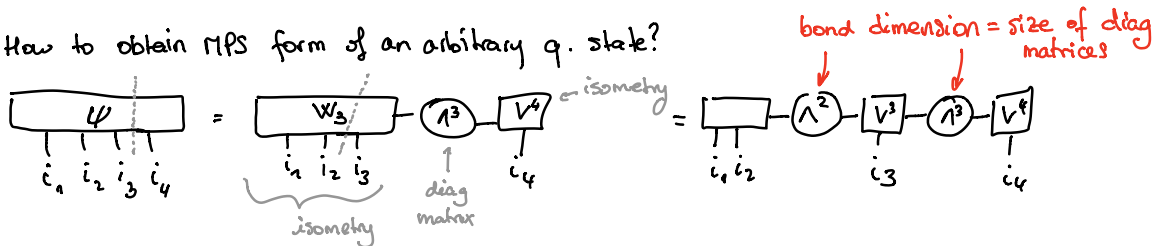
$$\text{here } A^{[i]} \in M_{D_i \times D'_{i+1}} \quad ; \quad D'_{i-1} = D_i \quad ; \quad D_1 = 1 \quad ; \quad D'_N = 1$$

$$\begin{array}{c} \boxed{\psi} \\ | \quad \dots \quad | \\ i_1 \quad \dots \quad i_N \end{array} = \begin{array}{c} \boxed{A^1} \quad \boxed{A^2} \quad \dots \quad \boxed{A^N} \\ | \quad \quad | \quad \quad \quad | \\ i_1 \quad i_2 \quad \quad \quad i_N \end{array} \quad D = \max D_i, D'_i = \text{bond dimension}$$

→ any state $|\psi\rangle$ can be written as a MPS with $D \sim d^{N/2}$ (exp bond dimension)

→ constant bond dimension $\Rightarrow$ area law since $S(S_{1,j}) = \log D_j \in D$ ← want fixed bond dimension

→ How to obtain MPS form of an arbitrary q. state?



→ iterating SVD's we get MPS form with bond dimension at step $i = \text{rank}_2 \rho_{i,i} = \# \text{supp } \Lambda^i$

→ how to obtain approx of $|\psi\rangle$ by MPS with small bond dimension?

Lemma 1 [MPS approximation with given bond dimension] [vc, lemma 1]

Given a state $\psi \in (\mathbb{C}^d)^{\otimes N}$, construct usual MPS form with matrices Λ^i (large bond dim)

Define approx $|\psi_D\rangle$ by truncating Λ^i to their largest D singular values. Then

$$\|\psi - \psi_D\|^2 \leq 2 \sum_{j=D+1}^{N-1} \epsilon_j(D) \quad \text{where} \quad \epsilon_j(D) = \sum_{i=D+1}^{N_j} \Lambda^j(i)$$

sing values forgotten when truncating

→ bound size of supp of Λ^i matrices $\leadsto S_0$ $\alpha=0$ Rényi entropy

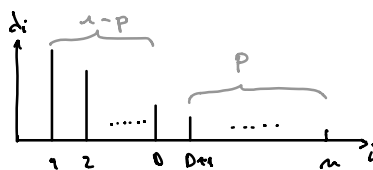
Lemma 2 let λ be a prob. vector (singular values of Λ^i). $\epsilon_D(\lambda) = \sum_{i=D+1} \lambda_i$ [vc, lemma 2]
 $\lambda_1 \geq \lambda_2 \geq \dots$

Then $\forall 0 \leq \alpha < 1$ $\log \epsilon_D(\lambda) \leq \frac{1-\alpha}{\alpha} \left[S_\alpha(\lambda) - \log \frac{D}{1-\alpha} \right]$
 want: ≤ 0

Proof idea

Consider:

$$\begin{aligned} \min & S_\alpha(\lambda) \\ \text{st.} & \sum_{i \geq D+1} \lambda_i = p \end{aligned}$$

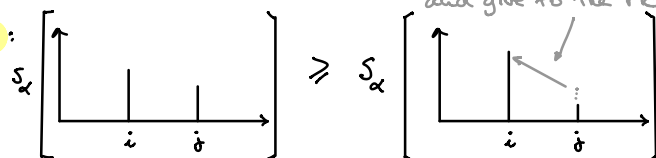


"take from the poor and give to the rich"

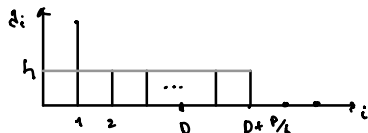
→ Rényi entropies are Schur-concave:

$$p \prec q \Rightarrow S_\alpha(p) \geq S_\alpha(q)$$

↑ majorization



→ under the "tail mass" constraint $\sum_{i \geq D+1} \lambda_i = p$, the optimizer must be of the form



$$\text{i.e. } \lambda_1 = 1 - p - (D-1)h; \lambda_2 = \dots = \lambda_{D+1} = h; \lambda_{D+2} = \dots = 0$$

↳ optimize over h to get inequality



Important remark

Hastings proves that the von Neumann entropy $S = S_{\alpha=1}$ is bounded, whereas the results above require $0 \leq \alpha < 1$. See [Has, section 2] for the reduction in this case.

→ more general results about approx. by MPS under entropy bound are given in [swvc]

[swvc]

$S_{\alpha} \sim$	const	$\log L$	$L^{\kappa} (\kappa < 1)$	L
$S_{\alpha < 1}$	approximable		undetermined	
$S \equiv S_1$				
$S_{\alpha > 1}$			inapproximable	

FIG. 1 (color online). Relation between scaling of block Rényi entropies and approximability by MPSs. In the “undetermined” region, nothing can be said about approximability just from looking at the scaling.

REFERENCES

[Has] Hastings - An area law for 1D q. syst (2007)

[vc] Verstraete, Cirac - MPS represent ground states faithfully

[swvc] Schuch, Wolf, Verstraete, Cirac - Entropy scaling and simulability by MPS