

① Classical and Q. Inf. Th.

Sets vs Hilbert spaces

classical Inf Th.

Finite alphabet

$$\Sigma_A = \{1, 2, \dots, d\}$$

different letters

$$i \neq j$$

Quantum Inf Th

fin dim, \mathbb{C} -Hilbert spaces

$$\mathcal{H}_A = \mathbb{C}^d =$$

$$\text{span} \{ |1\rangle, |2\rangle, \dots, |d\rangle \}$$

$$|x\rangle = \frac{1}{\sqrt{2}} (|1\rangle + |2\rangle)$$

superposition

orthogonal vectors

$$|x\rangle \perp |y\rangle$$

States

Classical

probability measures

supported on Σ_A

$$\mathbb{P}(\Sigma_A)$$

$$p \in \mathbb{R}^d \quad p_i \geq 0$$

$$\sum p_i = 1$$

p

extreme points

$$\{ \delta_i \}_{i \in \Sigma_A}$$

Quantum

density matrices

$$\rho \in \mathcal{B}(\mathcal{H}_A) \approx \text{M}_d^{\text{so}}(\mathbb{C})$$

ρ is positive semidefinite

$$\rho \geq 0$$

$$\text{Tr } \rho = 1 \quad \rho \in \mathcal{D}(\mathcal{H}_A)$$

$$\text{diag}(\rho)$$

$$\text{ext } \mathcal{D}(\mathcal{H}_A) = \text{rank-1}$$

$$\{ |x\rangle\langle x| : |x\rangle \text{ proj}, \|x\|=1, x \in \mathbb{C}^d \}$$

Example $d=2$ bits

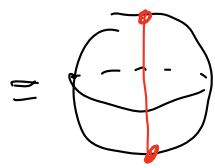
$$p = (1-t)\delta_0 + t\delta_1$$

$\mathbb{R}(\Sigma_2) = \text{segment}$



$d=2$ qubits

$\mathbb{D}(\mathbb{C}^2) = \text{Bloch ball}$



$$\rho = \frac{1}{2} (\mathbb{I} + x \cdot \sigma_x + y \cdot \sigma_y + z \cdot \sigma_z)$$

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \quad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$\rho \geq 0 \quad (\Rightarrow) \quad \|(x, y, z)\| \leq 1$$

center of the Bloch ball = $\mathbb{I}/2$

uniform distributions

$$p = \left(\frac{1}{d}, \frac{1}{d}, \dots, \frac{1}{d} \right)$$

$$\rho_* = \mathbb{I}_d / d$$

maximally mixed state

Bipartite and multipartite systems

Classical

Quantum

joint system $A+B$

\times
cartesian product

$$\Sigma_{AB} = \Sigma_A \times \Sigma_B$$

$$\left\{ (i, j) : \begin{array}{l} i \in \Sigma_A \\ j \in \Sigma_B \end{array} \right\}$$

\otimes
tensor product

$$\mathcal{H}_{AB} = \mathcal{H}_A \otimes \mathcal{H}_B$$

$$\mathbb{D}(\mathcal{H}_{AB}) \quad \rho_{AB} \geq 0$$

$$\text{Tr} \rho_{AB} = 1$$

$$P_{AB}(i,j) \geq 0$$

$$\sum_{ij} P_{AB}(i,j) = 1$$

product measures

$$P_{AB} = P_A \times P_B$$

$$P_{AB}(i,j) = P_A(i) P_B(j)$$

(independent r.v.)

$$\mathcal{H}_{AB} = \text{span} \{ |a\rangle \otimes |b\rangle : \begin{array}{l} a \in \mathcal{H}_A \\ b \in \mathcal{H}_B \end{array} \}$$

product states

$$\mathcal{P}_{AB} = \mathcal{P}_A \otimes \mathcal{P}_B$$

$$X \otimes Y = \begin{bmatrix} x_{11}y & x_{12}y & & \\ & & \ddots & \\ & & & x_{mm}y \end{bmatrix}$$

every bipartite state is a convex combination of product states.

not always the case!

there exist

entangled states

$$P_{AB} = \sum_{ij} P_{AB}(i,j) \cdot \delta_i \times \delta_j$$

convex comb. of prod. st.

$$\text{SEP}(A,B) :=$$

$$\text{Conv} \{ \mathcal{P}_A \otimes \mathcal{P}_B :$$

$$\mathcal{P}_A \in \mathcal{D}(\mathcal{H}_A)$$

$$\mathcal{P}_B \in \mathcal{D}(\mathcal{H}_B) \}$$

$$t_i \geq 0 \quad \sum t_i = 1$$

$$\mathcal{P}_{A,B}^{(i)} \in \mathcal{D}(\mathcal{H}_{A,B})$$

$$\text{SEP} \ni \rho = \sum t_i \rho_A^{(i)} \otimes \rho_B^{(i)}$$

extreme separable states

$$\text{ext}(\text{SEP}(A,B)) = \{ |ax\rangle \otimes |by\rangle : \begin{array}{l} x \in \mathcal{H}_A \\ y \in \mathcal{H}_B \end{array} \}$$

a pure state $z \in \mathcal{H}_{AB} = \mathcal{H}_A \otimes \mathcal{H}_B$

$$z \in \text{SEP} \Leftrightarrow z = x \otimes y$$

pure sep states in $\mathbb{C}^2 \otimes \mathbb{C}^2$:

$$\begin{aligned} z &= |1\rangle \otimes |1\rangle, \quad \frac{1}{\sqrt{2}} (|1\rangle \otimes |1\rangle + |1\rangle \otimes |2\rangle) \\ &= |1\rangle \otimes \frac{1}{\sqrt{2}} (|1\rangle + |2\rangle) \end{aligned}$$

entangled pure states.

$$z = \frac{1}{\sqrt{2}} (|1\rangle \otimes |1\rangle + |2\rangle \otimes |2\rangle) \neq |x\rangle \otimes |y\rangle$$

In general **maximally entangled state**

$$|\Omega\rangle = \frac{1}{\sqrt{d}} \sum_{i=1}^d |i\rangle \otimes |i\rangle \in \mathbb{C}^d \otimes \mathbb{C}^d$$

$$\omega = |\Omega\rangle\langle\Omega| \in \mathcal{D}(\mathbb{C}^d \otimes \mathbb{C}^d)$$

$$d=2 \quad \rho = \frac{1}{2} \begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix}$$

classical

marginalization

$$AB \rightarrow A$$

$$P_A(i) = \sum_j P_{AB}(i, j)$$

quantum

partial trace

$$\begin{aligned} \rho_{AB} &\rightarrow \rho_A \\ \rho_A &= [\text{id}_A \otimes \text{Tr}_B](\rho_{AB}) \\ &= \text{Tr}_B(\rho_{AB}) \end{aligned}$$

example $\text{Tr}_A \omega = \text{Tr}_B \omega = \text{Id}/d$

Quantum measurements

classical

$$\mathbb{P}(i) = p_A(i)$$

quantum $\rho \in \mathcal{D}(\mathcal{H}_A)$

POVM positive operator valued measure

$$M = \{M_x\}_{x \in \Sigma}$$

$$M_x \in \mathcal{B}(\mathcal{H}_A) \quad M_x \geq 0$$

$$\sum M_x = I_A$$

Born rule

$$\mathbb{P}(x) = \text{Tr}(M_x \cdot \rho)$$
$$\stackrel{\uparrow}{=} \langle M_x, \rho \rangle$$

Examples

- basis measurements

$$\Sigma = [d] = \{1, \dots, d\}$$

$$M_x = |x\rangle\langle x|$$

$$\mathbb{P}(x) = \langle x | \rho | x \rangle$$

- trivial measurements

$$M_x = q(x) \cdot I_A$$

$$\mathbb{P}(x) = q(x) \cdot \underbrace{\text{Tr}(I_A \cdot \rho)}_{=1}$$

State discrimination

Alice can prepare ρ_0 or ρ_1

w. prob p_0 p_1

$$p_0 + p_1 = 1$$

Bob receives ρ

has to decide

if $\rho = \rho_0$

or $\rho = \rho_1$

Win if Bob guesses correctly.

Want: $\mathbb{P}(\text{win}) = p_0 \mathbb{P}(\text{Bob answers 0} \mid \text{given } \rho = \rho_0) + p_1 \dots$

Bob can measure ρ using $M = (M_0, M_1)$

$$M_{0,1} \geq 0 \quad M_0 + M_1 = I_d.$$

$$\mathbb{P}(\text{win}) = p_0 \text{Tr}(M_0 \rho_0) + p_1 \text{Tr}(M_1 \rho_1)$$

Thm [Holevo, Helstrom]

$$p_0 = p_1 = 1/2$$

$$\max_{M_{0,1}} \mathbb{P}(\text{win}) = \frac{1}{2} + \frac{1}{4} \|\rho_0 - \rho_1\|_1$$

where $\|\cdot\|_1$ = trace norm
nuclear norm
Schatten-1 norm

$$\|X\|_1 = \text{Tr} |X| = \text{Tr} [(X^\dagger X)^{1/2}]$$

$$\|p - q\|_{TV} \Leftrightarrow \|\rho_p - \rho_q\|_1$$

$$\text{if } \rho_p = \text{diag}(p)$$

$$\rho_q = \text{diag}(q)$$

$$\sum_i |p(i) - q(i)|$$

$$\max \mathbb{P}(\text{win}) = \frac{1}{2} + \frac{1}{2} \|p - q\|_{TV}$$

sketch of proof. *diag* *diag*

$$\max_{\substack{M_0, M_1 \geq 0 \\ M_0 + M_1 = I}} p_0 \langle M_0, \rho_0 \rangle + p_1 \langle M_1, \rho_1 \rangle$$

$\frac{1}{2} \qquad \qquad \qquad \frac{1}{2}$

$$\max_{0 \leq M_0 \leq I} \frac{1}{2} \langle M_0, \rho_0 - \rho_1 \rangle + \underbrace{\frac{1}{2} \langle I, \rho_1 \rangle}_{= 1/2}$$

$$\text{arg max} = \mathbb{1}(\rho_0 - \rho_1)_+$$

$$\text{if } \rho_0 - \rho_1 = \sum \lambda_i |x_i\rangle\langle x_i|$$

$$M = \sum \mathbb{1}_{\lambda_i \geq 0} |x_i\rangle\langle x_i|$$

$$f_{\max} = \frac{1}{2} + \frac{1}{2} \sum \lambda_i \cdot \mathbb{1}_{\lambda_i \geq 0}$$

$$|\rho_0 - \rho_1| = \sum |\lambda_i| |x_i\rangle\langle x_i|$$

$$\begin{aligned} \|\rho_0 - \rho_1\|_1 &= \text{Tr} |\rho_0 - \rho_1| = \sum |\lambda_i| = \sum_{\lambda_i \geq 0} \lambda_i + \sum_{\lambda_i < 0} (-\lambda_i) \\ &= 2 \sum_{\lambda_i \geq 0} \lambda_i \end{aligned}$$

$$\text{Since } \text{Tr}(\rho_0 - \rho_1) = 0 = \sum \lambda_i$$

② Quantum Channels

Classical

cl. channel: transformations of classical states i.e. prob. meas.

$$N: \mathbb{R}^{|\Sigma_A|} \rightarrow \mathbb{R}^{|\Sigma_B|}$$

$$\forall a \in \Sigma_a$$

$$b \in \Sigma_b$$

$$N_{b|a} = \langle \delta_b | N(\delta_a) \rangle \geq 0 \rightarrow$$



$$\forall a \quad \sum_b N_{b|a} = 1$$

Markov kernels

Conditional prob., ...

Stochastic mat.

Why complete positivity?

(automatic in the classical case)

\exists positive maps which are not completely positive e.g. the transposition

$$\text{transp}: M_2 \rightarrow M_2 \quad \text{is positive}$$

$$X \mapsto X^T$$

$$(\text{transp} \otimes \text{id}_2)(\omega_2) \not\geq 0$$

Quantum

q. channels = transformations of q. states (density matrices)

$$\Phi: \mathcal{B}(\mathcal{H}_A) \rightarrow \mathcal{B}(\mathcal{H}_B)$$

linear

Φ is completely positive ^{CP}

$\forall \mathcal{H}_R \quad \Phi \otimes \text{id}_R$ positive

$$\forall \mathbb{Z} \geq 0 \quad (\Phi \otimes \text{id}_R)(\mathbb{Z}) \geq 0$$

Φ is trace preserving TP

$$\forall X \quad \text{Tr} \Phi(X) = \text{Tr} X$$

Choi matrix

$$\Phi: \mathcal{B}(\mathcal{H}_A) \rightarrow \mathcal{B}(\mathcal{H}_B) \text{ linear}$$
$$\rightarrow C_\Phi := \sum_{i,j=1}^{\dim \mathcal{H}_A} |i\rangle\langle j| \otimes \Phi(|i\rangle\langle j|)$$

\Rightarrow

$$\mathcal{B}(\mathcal{H}_A \otimes \mathcal{H}_B) \quad E_{ij}$$

Choi's theorem

$$\Phi \text{ is CP} \Leftrightarrow C_\Phi \geq 0$$

$$\Phi \text{ is TP} \Leftrightarrow \text{Tr}_B C_\Phi = I_A$$

Examples

Classical

$$N_{b|a} = \frac{1}{|\Sigma_B|}$$

$$N_{b|a} = \mathbb{1}_{b=\pi(a)}$$
$$\Sigma_A = \Sigma_B, \pi \in S_\Sigma$$

permutation

Quantum

depolarizing channel

$$\Delta(X) = \text{Tr}(X) \frac{I_B}{\dim \mathcal{H}_B}$$

unitary channels

$$\Phi_U(X) = U X U^\dagger$$

for $U \in \mathcal{U}_{\mathcal{H}_A = \mathcal{H}_B}$

measurement channels

$$M = (M_x)_{x=1}^k \text{ a POVM}$$

$$(M_x \geq 0, \sum_x M_x = I)$$

$$\Phi_M(\rho) = \sum_{x=1}^k \langle M_x, \rho \rangle \cdot |x\rangle\langle x|$$

Kraus decomposition Φ is CP & TP \Leftrightarrow

$$\exists K_1, \dots, K_r : \mathcal{H}_A \rightarrow \mathcal{H}_B \text{ s.t.}$$
$$\Phi(X) = \sum_{i=1}^r K_i X K_i^*$$
$$\sum_{i=1}^r K_i^* K_i = I_A$$

Stinespring dilation Φ is CP & TP \Leftrightarrow

$$\exists V : \mathcal{H}_A \rightarrow \mathcal{H}_B \otimes \mathcal{H}_E \text{ isometry}$$

s.t. $\Phi(X) = \text{Tr}_E(V X V^*)$

Kraus
↓
Stinespring

$$Vx = \sum_{i=1}^r K_i x \otimes |i\rangle$$

$$\mathcal{H}_E = \text{span} \{ |1\rangle, |2\rangle, \dots, |r\rangle \}$$

$$\min r = \text{rank } C_\Phi$$

N Markov matrix $\rightsquigarrow \Phi^N$

$$\Phi^N(X) = \sum_{\substack{a \in \Sigma_A \\ b \in \Sigma_B}} N_{ba} (a|X|a) |b\rangle\langle b|$$

$$C_{\Phi^N} = \sum N_{ba} |a\rangle\langle a| \otimes |b\rangle\langle b|$$

Different notions of positivity for quantum maps

Map $\Phi: \mathcal{B}(\mathcal{H}_A) \rightarrow \mathcal{B}(\mathcal{H}_B)$

positive maps

$$X \geq 0 \Rightarrow \Phi(X) \geq 0$$

Choi matrix $C_\Phi \in \mathcal{B}(\mathcal{H}_A \otimes \mathcal{H}_B)$

block-positive matrix

$$\langle x \otimes y | C_\Phi | x \otimes y \rangle \geq 0$$

$$\forall x \in \mathcal{H}_A \quad \forall y \in \mathcal{H}_B$$

completely positive maps

positive semidef.

$$\forall \mathcal{H}_R \quad Z_{AR} \geq 0 \Rightarrow [\Phi \otimes \text{id}_R](Z) \geq 0$$

$$\forall z \in \mathcal{H}_A \otimes \mathcal{H}_B$$

$$\langle z | C_\Phi | z \rangle \geq 0$$

dual

G

entanglement breaking

(super-positive)

$$\forall \mathcal{H}_R \quad Z_{AR} \geq 0 \Rightarrow [\Phi \otimes \text{id}_R](Z) \in \text{SEP}(B, R)$$

$$C_\Phi \in \text{SEP}(A, B)$$

③ Classical capacity of quantum channels

Classical Shannon

$$H(p) = - \sum_{i \in \Sigma} p_i \log p_i$$

$$H(p) \in [0, \log |\Sigma|]$$

$$p = \delta_i$$

$$p = \left(\frac{1}{|\Sigma|}, \dots, \frac{1}{|\Sigma|} \right)$$

$$H(A)_p$$

Quantum von Neumann

$$S(\rho) = - \text{Tr} \rho \log \rho$$

$$= H(\lambda_\rho)$$

↑ spectrum of ρ

$$S(\rho) \in [0, \log \dim \mathcal{H}]$$

$$\rho = |x\rangle\langle x|$$

pure state

$$\rho = \frac{I}{\dim \mathcal{H}}$$

Conditional entropy

$$H(A|B)_p = H(AB)_p - H(B)_p \geq 0$$

$$\begin{aligned} S(A|B)_\rho &= S(AB)_\rho - S(B)_\rho \\ &= S(\rho_{AB}) - S(\rho_B) \end{aligned}$$

can be negative!

$$\begin{aligned} \rho_{AB} &= \text{max. ent. state} = \omega \\ &= |\Omega\rangle\langle\Omega| \end{aligned}$$

$$|\Omega\rangle = \frac{1}{d} \sum |i\rangle \otimes |i\rangle$$

$$S(\rho_{AB}) = 0$$

$$S(\rho_B = \frac{I}{d}) = \log d$$

relative entropy

$$D_{KL}(p \parallel q) = \sum p_i \log \frac{p_i}{q_i} \quad D(\rho \parallel \sigma) = \text{Tr} \rho (\log \rho - \log \sigma)$$

$$D(\rho \parallel \sigma) \geq 0$$

Data-processing inequality

$\forall \Phi$ q. channel

$$D(\rho \parallel \sigma) \geq D(\Phi(\rho) \parallel \Phi(\sigma))$$

Mutual information

$$I(A:B)_p = H(A)_p + H(B)_p - H(AB)_p = H(B)_p - H(B|A)_p$$

$$I(A:B)_\rho = S(A)_\rho + S(B)_\rho - S(AB)_\rho$$

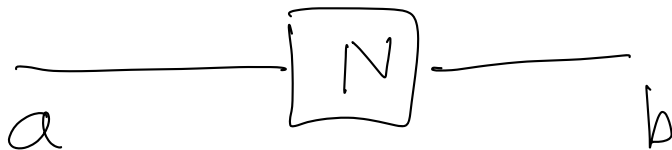
$$I(A:B) \geq 0 \quad = 0 \quad \text{iff} \quad P_{AB} = P_A \times P_B$$

$$\rho_{AB} = \rho_A \otimes \rho_B$$

Shannon's channel coding theorem

N is a classical channel

$$N = (N_{b|a})$$



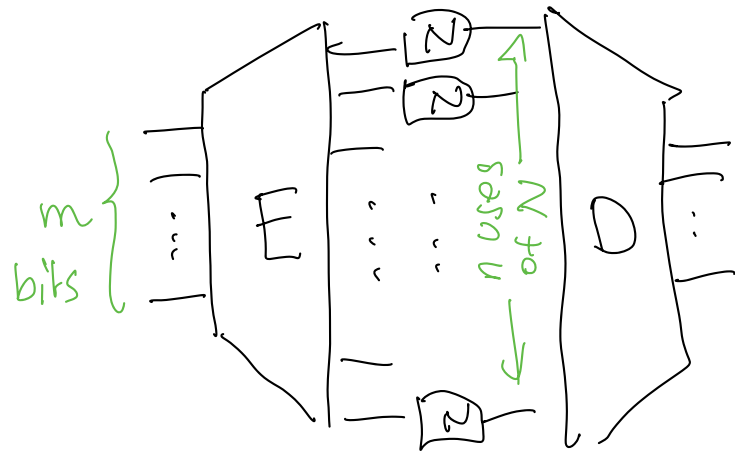
$$N: \sum_A \rightarrow \sum_B$$

coding scheme (m, n, δ) for N is.

a pair $E: \{0,1\}^m \rightarrow \sum_A^n$

$D: \sum_B^m \rightarrow \{0,1\}^m$ such that

$$\forall i \in \{0,1\}^m \quad \mathbb{P}(D \circ N \circ E(i) = i) \geq 1 - \delta$$



achievable rate R

R is achievable if $\forall n \exists (m_n, n, \delta_n)$ coding scheme for N s.t. $\frac{m_n}{n} \rightarrow R$ $\delta_n \rightarrow 0$

Capacity $C(N) = \sup \{ R : R \text{ is achievable} \}$

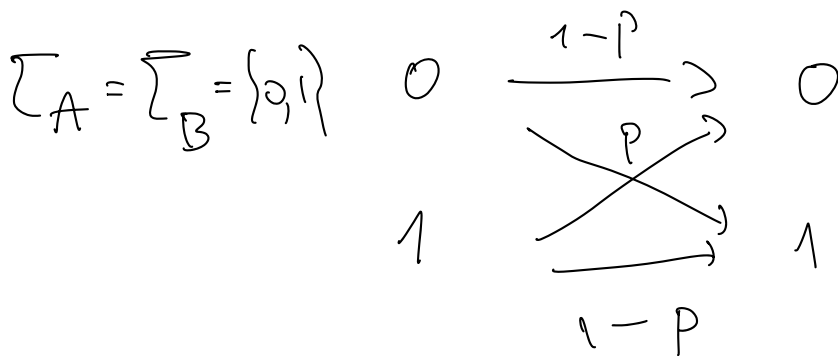
Theorem

$$C(N) = \sup_{P_A} I(A : B)_{P_{AB}^N}$$

$$P_{AB}^N(a, b) = P_A(a) N_{b|a}$$

Example

Binary symmetric channel.



BSC_p

$$C(\text{BSC}_p) = 1 - H(p, 1-p)$$

Classical capacity of a quantum channel

coding scheme for $\Phi : \mathcal{B}(\mathcal{H}_A) \rightarrow \mathcal{B}(\mathcal{H}_B)$
(m, n, δ)

$$E: \{0,1\}^m \rightarrow \mathcal{B}(\mathcal{H}_A^{\otimes n})$$

$$D: \mathcal{B}(\mathcal{H}_B^{\otimes n}) \rightarrow \{0,1\}^m \text{ POVM}$$

$$\forall i \in \{0,1\}^m \quad \langle D_i, \overline{\Phi}^{\otimes n}[E(i)] \rangle \geq 1 - \delta$$

achievable rate: $\forall n \exists (m_n, n, \delta_n)$ coding st

$$\frac{m_n}{n} \rightarrow R, \quad \delta_n \rightarrow 0$$

capacity $C(\overline{\Phi}) = \max \{R: R \text{ is achievable}\}$

$$D: \mathcal{B}(\mathcal{H}_B^{\otimes n}) \rightarrow \{0,1\}^m \text{ decoding channel.}$$

$\approx \text{diag } \mathcal{M}_{2^m}$

$$D(\rho) = \sum_{i \in \{0,1\}^m} \langle D_i, \rho \rangle \cdot |i\rangle\langle i|$$

Theorem [HSW]

$$C(\overline{\Phi}) = \lim_{n \rightarrow \infty} \frac{1}{n} \chi(\overline{\Phi}^{\otimes n})$$

where

$$\chi(\Psi) = \max_{\{p(i), \rho_i\}_{i \in \Sigma}} S\left(\sum_i p_i \Psi(\rho_i)\right) - \sum_i p_i S(\Psi(\rho_i))$$

Holevo capacity of Ψ

$$- \sum_i p_i S(\Psi(\rho_i))$$

if χ were additive i.e.

$$\chi(\Phi_1 \otimes \Phi_2) = \chi(\Phi_1) + \chi(\Phi_2)$$

then $C(\Phi) = \chi(\Phi)$ **No!**

Fact [Shor]

χ is additive $\Leftrightarrow S_{\min}$ additive

$$S_{\min}(\Phi) = \min_{\rho \in \mathcal{D}(\mathcal{H}_A)} S(\Phi(\rho)) = \min_{\sigma \in \Phi(\mathcal{D}(\mathcal{H}_A))} S(\sigma)$$

$$\chi(\Phi) = \max_{\rho, \gamma} \underbrace{S(\sum p_i \Phi(\rho_i))}_{\leq \log \dim \mathcal{H}_B} - \underbrace{\sum p_i S(\Phi(\rho_i))}_{\geq S_{\min}}$$

$$\leq \log \dim \mathcal{H}_B - S_{\min}(\Phi)$$

• if Φ is covariant \exists irreps U_g, V_g

$$\forall x \quad V_g \Phi(x) V_g^\dagger = \Phi(U_g x U_g^\dagger)$$

then $\chi(\Phi) = \log \dim \mathcal{H}_B - S_{\min}(\Phi)$

Example depolarizing channel

$$\Delta_\lambda(\rho) = (1-\lambda)\rho + \lambda \frac{I}{d}$$

$$\Delta_\lambda: \mathcal{M}_d \rightarrow \mathcal{M}_d \quad \mathcal{H}_A = \mathcal{H}_B = \mathbb{C}^d$$

Fact [King] $\forall \Phi$

$$\chi(\Delta_\lambda \otimes \Phi) = \chi(\Delta_\lambda) + \chi(\Phi)$$

$$C(\Delta_\lambda) = \chi(\Delta_\lambda) = \log d - S_{\min}(\Delta_\lambda)$$

$$= \log d - H\left(1-\lambda + \frac{\lambda}{d}, \frac{\lambda}{d}, \dots, \frac{\lambda}{d}\right)$$

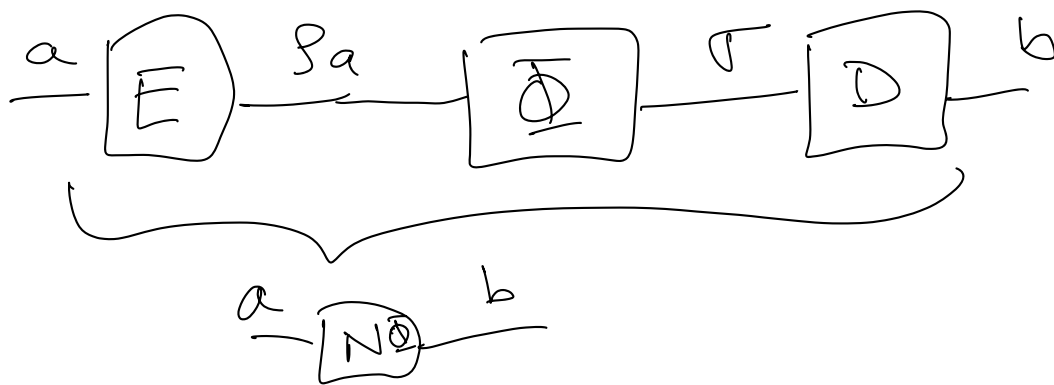
d-1 times

Relation between Shannon's and HSW theorems

$$\Phi: A \rightarrow B$$

$$N_{b|a}^\Phi = \langle M_b, \Phi(p_a) \rangle$$

Here E prepares p_a with prob p_a
 D measures (M_b)



$$C(N^\Phi) = \max_{P_A} I(A:B)_{P_{AB}^\Phi}$$

$$\max_{\rho_a, M_b} C(N \Phi) = \max_{\rho_a, \rho_a, M_b} I(A:B)_{P_{AB}}$$

$$P_{AB}(a, b) = P_A(a) \langle M_b, \Phi(\rho_a) \rangle$$

$$= \max_{\rho_a, \rho_a} \max_{M_b} I(A:B)_{P_{AB}}$$

$$\underbrace{\hspace{15em}}_{I_{\text{acc}}(\mathcal{E}_{\rho_a(a), \rho_a})}$$

accessible information

\mathcal{E} = state ensemble

$$\rho_A(a), \Phi(\rho_a)$$

$$I_{\text{acc}}(\mathcal{E}) = \max_{M_b} I(A:B)_{P_{AB}}$$

$$\leq \chi(\mathcal{E})$$

↑

Holevo

$$\chi(\mathcal{E}) = S\left(\sum_A P_A(a) \Phi(\rho_a)\right) -$$

$$- \sum_a P_A(a) S(\Phi(\rho_a))$$